Simple Edit Distance

It's easy in principle to compute edit distance with finite state transducers. We construct a transducer that takes each symbol in the alphabet to itself with cost 0, and takes each symbol to a different symbol, or to/from epsilon with cost 1. This transducer is called a *flower transducer* because of its appearance.

We then compose an FST corresponding to the first string with this transducer, compose the result with an FST corresponding to the second string, and compute the cost of the shortest path.
In [4]:
```
flower = make_flower("AB")
show_fst(flower)
```

In [5]:
```
fst1 = FST()
fst1.AddString("AABAAAA")
fst2 = FST()
fst2.AddString("AABBABAB")
```

In [6]:
```
templ = FST()
onfst.ArcSortOutput(fst1)
onfst.ArcSortInput(flower)
onfst.Compose(fst1,flower,templ)
show_fst(templ)
```
Factoring the Edit Distance Transducer

The problem with the previous transducer is that it gets very large very quickly when composed with the original string. In fact, the size ends up being quadratic.

We can fix this by introducing some additional symbols. (Here, we're just using ASCII symbols to represent insertion, deletion, and substitution, but we could be using something fancier.)
epsilon = 0
insertion = ord("#")
deletion = ord("_")
substitution = ord("~")

def make_left(chars):
    fst = FST()
    s = fst.AddState()
    fst.SetStart(s)
    fst.SetFinal(s, 0.0)
    for c in chars:
        c = ord(c)
        fst.AddArc(s, c, c, 0.0, s)
        fst.AddArc(s, c, substitution, 0.5, s)
        fst.AddArc(s, c, deletion, 0.5, s)
    return fst

def make_right(chars):
    fst = FST()
    s = fst.AddState()
    fst.SetStart(s)
    fst.SetFinal(s, 0.0)
    for c in chars:
        c = ord(c)
        fst.AddArc(s, c, c, 0.0, s)
        fst.AddArc(s, substitution, c, 0.5, s)
        fst.AddArc(s, insertion, c, 0.5, s)
    return fst

temp1 = FST()
temp2 = FST()
openfst.Compose(fst1, make_left("AB"), temp1)
openfst.Compose(make_right("AB"), fst2, temp2)
print fsts size(temp1), fsts size(temp2)

(8, 29) (9, 33)
Limited Contiguous Insertions / Deletions
A second way in which we can make edit distance computations more efficient is to limit the number of consecutive deletions/insertions that can occur.

(Think about what constraint this corresponds to for a "manual" computation of the edit distance.)

```
In [19]: epsilon = 0  
def make_edit1(chars):
    fst = FST()
    s = fst.AddState()
    s2 = fst.AddState()
    fst.SetStart(s)
    fst.SetFinal(s, 0.0)
    fst.SetFinal(s2, 0.0)
    for c in chars:
        c = ord(c)
        fst.AddArc(s, c, c, 0.0, s)
        fst.AddArc(s, c, epsilon, 1.0, s2)
        fst.AddArc(s2, c, c, 1.0, s2)
    for c2 in chars:
        c2 = ord(c2)
        fst.AddArc(s, c, c2, 1.0, s)
        fst.AddArc(s2, c, c2, 1.0, s)
    return fst
```
In [20]:
    temp1 = FST()
    openfst.ArcSortOutput(fst1)
    efst = make_edit1("AB")
    openfst.ArcSortInput(efst)
    openfst.Compose(fst1, efst, temp1)
    show_fst(temp1)
    temp2 = FST()
    openfst.ArcSortOutput(temp1)
    openfst.ArcSortInput(fst2)
    openfst.Compose(temp1, fst2, temp2)
    show_fst(temp2)
    print ftsize(temp2)

(56, 115)

In [21]:
    result = FST()
    openfst.ShortestPath(temp2, result, 1)
    show_fst(result)
Oracle Edit Distance

The regular edit distance is limited to computing the best match between two strings. However, with finite state transducers, we can compute the best match between two sets of strings.

In fact, the strings can even have weights or costs associated with them.

```python
In [22]:
    temp1 = FST()
    openfst.ArcSortOutput(fst1)
    efst = make_flower("AB")
    openfst.ArcSortInput(efst)
    openfst.Compose(fst1,efst,temp1)
    show_fst(temp1)
    temp2 = FST()
    openfst.ArcSortOutput(temp1)
    openfst.ArcSortInput(fst2)
    openfst.Compose(temp1,fst2,temp2)
    show_fst(temp2)
    print fstsize(temp2)

(72, 211)
```

```python
In [23]:
    result = FST()
    openfst.ShortestPath(temp2,result,1)
    show_fst(result)
```
Let's actually use an English dictionary (albeit a small one) to find the closest string in this set of hypotheses to some dictionary word.

Often, we'd simply precompute such a minimized representation of the dictionary and then use it for further recognition.

```
In [34]: # recognition output
fst1 = FST()
fst1.AddString("quck", 1.1)
fst1.AddString("qwck", 1.3)
fst1.AddString("quidc", 2.3)
fst1 = minimize(fst1)
show_fst(fst1)
```

Let's actually use an English dictionary (albeit a small one) to find the closest string in this set of hypotheses to some dictionary word.

```
In [35]: # English dictionary
fst2 = FST()
with open("basic-english.txt") as stream:
    for line in stream.readlines():
        line = line.strip()
        fst2.AddString(line)
print fstsize(fst2)
fst2 = minimize(fst2)
print fstsize(fst2)
```

```
(4459, 4458)
(746, 1541)
```

```
In [36]: temp2 = FST()
onefst.ArcSortOutput(ascii_right)
onefst.ArcSortInput(fst2)
onefst.Compose(ascii_right,fst2,temp2)
print fstsize(temp2)
```

```
(746, 5369)
```

Note that might actually try to "minimize" this dictionary transducer. However, this doesn't work because this is a "non-functional FST". In order to minimize this transducer, one needs to encode the labels, determinize, then decode, and finally minimize.

```
In [37]: # fst2 = minimize(fst2)
```

Now let's compose the set of hypotheses with the dictionary.
In [38]:
    temp1 = FST()
    openfst.Compose(fst1, ascii_left, temp1)
    print fstsize(temp1), fstsize(temp2)
    temp3 = FST()
    openfst.ArcSortOutput(temp1)
    openfst.ArcSortInput(temp2)
    openfst.Compose(temp1, temp2, temp3)
    print fstsize(temp3)
    result = FST()
    openfst.ShortestPath(temp3, result, 1)
    show_fst(result)

(8, 35) (746, 5369)
(5968, 33447)

In [27]: